Parallel Programming and Optimization with GCC

Diego Novillo
dnovillo@google.com
• Parallelism models

• Architectural overview

• Parallelism features in GCC

• Optimizing large programs
  – Whole program mode
  – Profile guided optimizations
Use hardware concurrency for increased
– Performance
– Problem size

Two main models
– Shared memory
– Distributed memory

Nature of problem dictates
– Computation/communication ratio
– Hardware requirements
- Processors share common memory
- Implicit communication
- Explicit synchronization
- Simple to program but hidden side-effects
Distributed Memory

- Each processor has its own private memory
- Explicit communication
- Explicit synchronization
- Difficult to program but no/few hidden side-effects
GCC supports four concurrency models:

- **Easy**
  - ILP: automatic, no user control, not intrusive
  - Vectorization: automatic, compiler option, not intrusive
  - OpenMP: manual, compiler directives, somewhat intrusive
  - MPI: manual, special libraries, very intrusive

Ease of use not necessarily related to speedups!
Vectorization

Perform multiple array computations at once

Two distinct phases

- Analysis → high-level
- Transformation → low-level

Successful analysis depends on

- Data dependency analysis
- Alias analysis
- Pattern matching

Suitable only on loop intensive code
Vectorization

int a[256], b[256], c[256];

foo ()
{
    for (i = 0; i < 256; i++)
        a[i] = b[i] + c[i];
}

Vectorized

.L2:
    movdqa  c(%eax), %xmm0
    paddb   b(%eax), %xmm0
    movdqa  %xmm0, a(%eax)
    addl    $16, %eax
    cmpl    $1024, %eax
    jne     .L2

Scalar

.L2:
    movl    c(,%edx,4), %eax
    addl    b(,%edx,4), %eax
    movl    %eax, a(,%edx,4)
    addl    $1, %edx
    cmpl    $256, %edx
    jne     .L2

(~2x on P4)
OpenMP – Programming Model

Based on fork/join semantics
– Master thread spawns teams of children threads
– All threads share common memory

Allows sequential and parallel execution
```c
#include <omp.h>

main()
{
    #pragma omp parallel
    printf ("[%d] Hello\n", omp_get_thread_num());
}
```

$ gcc -fopenmp -o hello hello.c
$ ./hello
[2] Hello
[3] Hello
[0] Hello  ← Master thread

$ gcc -o hello hello.c
$ ./hello
[0] Hello
## Optimization Options

<table>
<thead>
<tr>
<th>Level</th>
<th>Transformations</th>
<th>Speed</th>
<th>Debuggability</th>
</tr>
</thead>
<tbody>
<tr>
<td>-O0</td>
<td>None (default)</td>
<td>Slow</td>
<td>Very good</td>
</tr>
<tr>
<td>-O1</td>
<td>Few</td>
<td>Not so fast</td>
<td>Good</td>
</tr>
<tr>
<td>-O2</td>
<td>Many</td>
<td>Fast</td>
<td>Poor</td>
</tr>
<tr>
<td>-Os</td>
<td>Same as -O2 + size</td>
<td>N/A</td>
<td>Poor</td>
</tr>
<tr>
<td>-O3</td>
<td>Most</td>
<td>Faster</td>
<td>Very poor</td>
</tr>
<tr>
<td>-O4</td>
<td>Nothing beyond -O3</td>
<td>N/A</td>
<td>N/A</td>
</tr>
</tbody>
</table>

- It may be faster than -O2 due to smaller footprint.
Optimization Options

Optimizations done at two levels

- Target independent, controlled with -f
- Target dependent, controlled with -m

There are more than 100 passes

Not all can be controlled with -f/-m

-Ox is **not** equivalent to a bunch of -f/-m

Use `-fverbose-asm -save-temps` to determine what flags were enabled by -Ox

Use `-fno-...` to disable a specific pass
Enabling additional optimizations

Not every available optimization is enabled by -Ox

- ftree-vectorize
- ftree-loop-linear
- ftree-loop-im
- funswitch-loops (-O3)
- funroll-loops
- finline-functions (-O3)
- ffast-math

Hundreds of -f and -m flags in the documentation
Optimizations are limited by the amount of code that the compiler can see at once. Current technology only works across one file at a time. Compiler must be able to work across file boundaries.
Optimizing Very Large Programs

Problem
Thousands of files, millions of functions, tens of gigabytes Massivememory/computation complexity for a single machine
WHOPR Architecture - 1

Front End
- Source
- LGEN
- Local Generation (LGEN)

Middle End
- Whole Program Analysis (WPA)
- LTRANS
- Local Transformation (LTRANS)

Back End
- Assembler
- LTRANS
- Assembler

Generate GIMPLE (parallel)
Make global optimization decisions (sequential)
Apply global decisions locally (parallel)
Compilation proceeds in 3 main phases:

- **LGEN (Local GENeration)**
  - Writes out GIMPLE
  - Produces summary information

- **WPA (Whole Program Analysis)**
  - Reads summary information
  - Aggregates local callgraphs into global callgraph
  - Produces global optimization plan

- **LTRANS (Local TRANSformation)**
  - Applies global optimization plan to individual files
  - Performs intra-procedural optimizations
  - Generates final code
• Phases 1 (LGEN) and 3 (LTRANS) are massively parallel

• Phase 2 (WPA) is fan-in/fan-out serialization point
  ○ Only operates with call graph and symbols
  ○ Transitive closure analysis not computationally expensive

• Scalability provided by splitting analysis and final code generation
  ○ Restricts types of applicable optimizations
  ○ For smaller applications, LTRANS provides full IPA functionality (whole program in memory)
Profile Guided Optimization

• Three phases
  – Profile code generation: Compile with `-fprofile-generate`
  – Training run: Run code as usual
  – Feedback optimization: Recompile with `-fprofile-use`

• Allows very aggressive optimizations based on accurate cost models
  – Provided that training run is representative!

• Compilation process significantly more expensive

• May not be applicable in all cases
GProf

Probes inserted automatically by compiler

Compile and link application with -pg

Run application as usual

Use gprof to analyze output file gmon.out

$ gcc -pg -O2 -o matmul matmul.c
$ ./matmul
$ gprof ./matmul
OProfile

- System-wide profiler.
- No modifications to source code
- Samples hardware counters to collect profiling information
- User specifies which hardware counter to sample
- Needs super-user access to start
  - Start Oprofiler daemon
  - Run application
- Use reporting program to read collected profile
Profile Guided Optimization Advances

• Instrument → Run → Recompile cycle too demanding
• New feature being developed to use hardware counters
  1. Program compiled as usual
  2. Runs in production environment with hardware counters enabled
  3. Subsequent recompilations use profile information from hardware counters

This allows for always-on, transparent profile feedback
• There is no “right” choice
• Granularity of work main indicator
• Evaluate complexity ↔ speedup trade-offs
• Combined approach for complex applications
• Algorithms matter!
• Good sequential algorithms may make bad parallel ones
Conclusions

• Performance tuning goes beyond random compiler flags
• Profiling tools are important to study behaviour
• Each tool is best suited for a specific usage
  – Try different flags and use /usr/bin/time to measure
  – Oprofile → system wide
  – Gprof → intrusive but useful to isolate profiling scope
  – Compiler dumps to determine source of problem
• New advances in instrumentation and whole program compilation will simplify things