Parallel Programming and Optimization with GCC

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- Parallelism models
- Architectural overview
- Parallelism features in GCC
- Optimizing large programs
 - Whole program mode
 - Profile guided optimizations

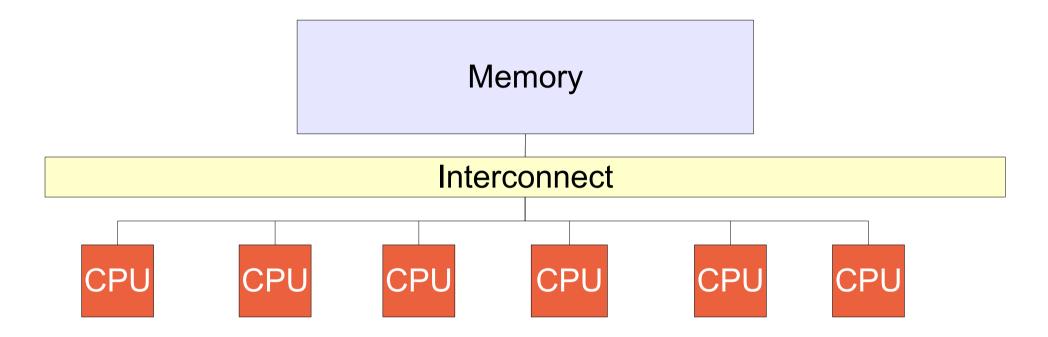
Parallel Computing



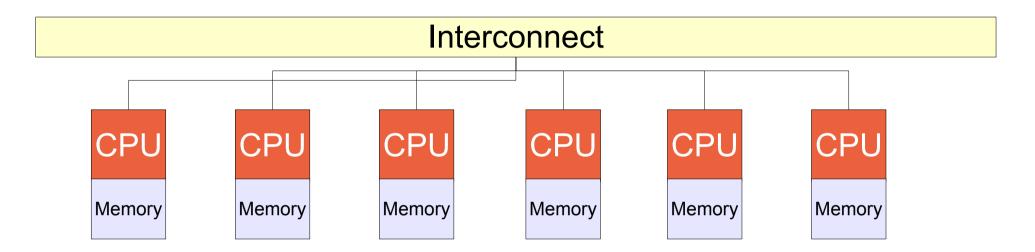
Use hardware concurrency for increased

- Performance
- Problem size
- Two main models
 - Shared memory
 - Distributed memory
- Nature of problem dictates
 - Computation/communication ratio
 - Hardware requirements





- Processors share common memory
- Implicit communication
- Explicit synchronization
- Simple to program but hidden side-effects

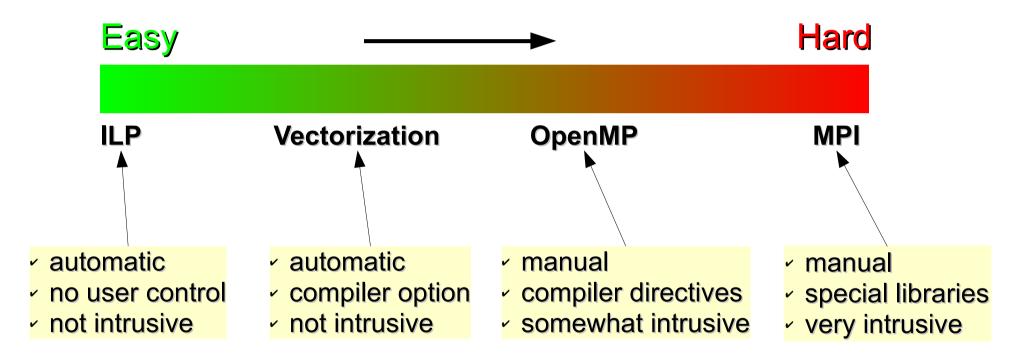


- Each processor has its own private memory
- Explicit communication
- Explicit synchronization
- Difficult to program but no/few hidden side-effects

Goog



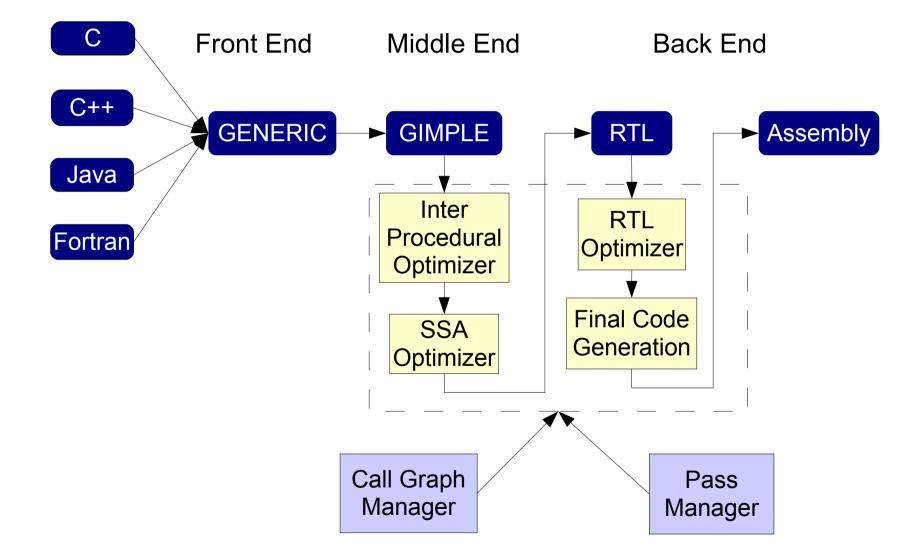
GCC supports four concurrency models



Ease of use not necessarily related to speedups!

GCC Architecture







Perform multiple array computations at once

Two distinct phases

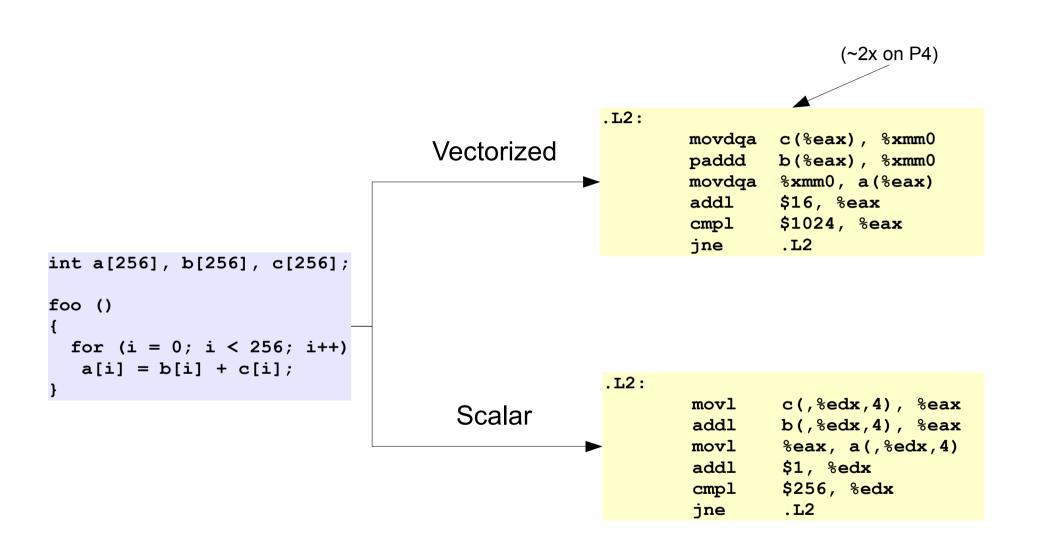
- -Analysis \rightarrow high-level
- $-Transformation \rightarrow low-level$

Successful analysis depends on

- -Data dependency analysis
- -Alias analysis
- -Pattern matching

Suitable only on loop intensive code

Vectorization

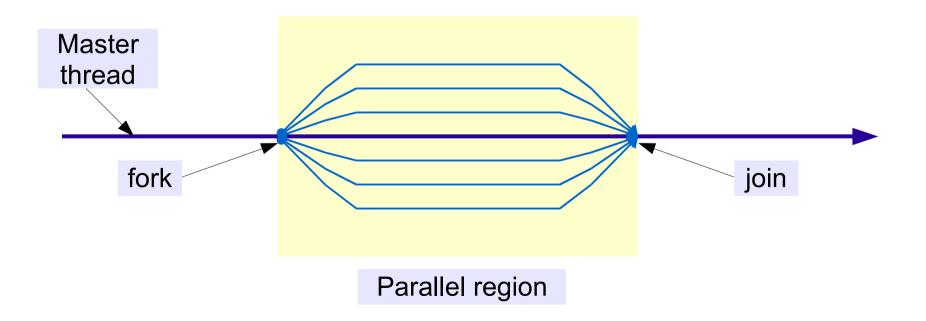




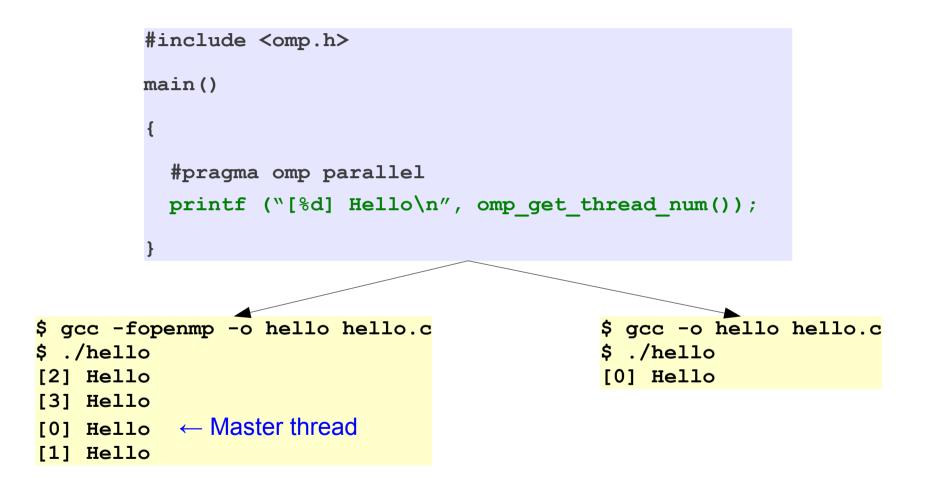
Based on fork/join semantics

- Master thread spawns teams of children threads
- All threads share common memory

Allows sequential and parallel execution







Optimization Options

Level	Transformations	Speed	Debuggability
-00	None (default)	Slow	Very good
-01	Few	Not so fast	Good
-02	Many	Fast	Poor
-Os	Same as -O2 + size	N/A	Poor
-03	Most	Faster	Very poor
-04	Nothing beyond -O3	N/A	N/A
It may be faster than -O2 due to smaller footprint			



Optimizations done at two levels

- Target independent, controlled with -f
- Target dependent, controlled with -m

There are more than 100 passes

Not all can be controlled with -f/-m

-Ox is **not** equivalent to a bunch of -f/-m

Use -fverbose-asm -save-temps to determine what flags were enabled by -Ox

Use -fno-... to disable a specific pass

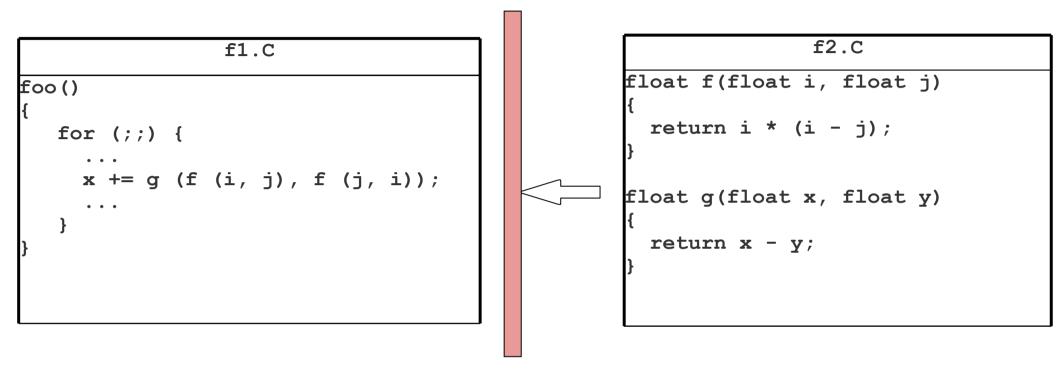
Not every available optimization is enabled by -Ox

- -ftree-vectorize
- -ftree-loop-linear
- -ftree-loop-im
- -funswitch-loops (-03)
- -funroll-loops
- -finline-functions (-03)
- -ffast-math

Hundreds of -f and -m flags in the documentation

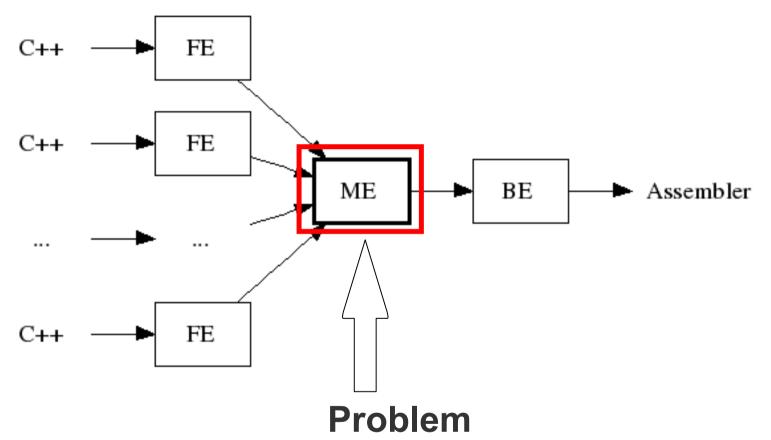
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Optimizing Very Large Programs Google



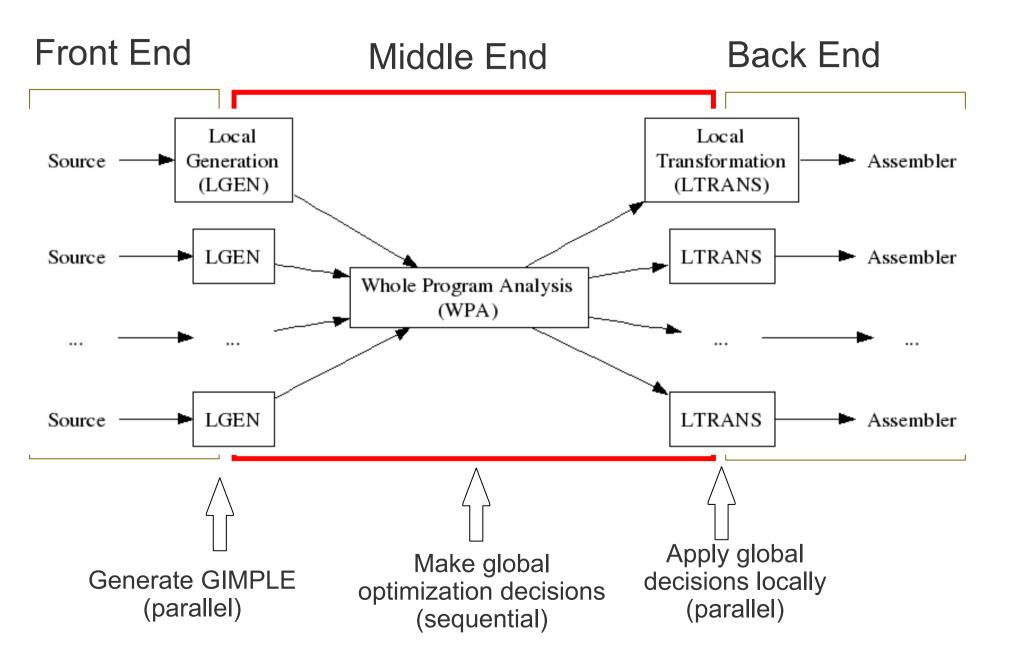
Optimizations are limited by the amount of code that the compiler can see at once Current technology only works across one file at a time Compiler must be able to work across file boundaries

Optimizing Very Large Programs Google



Thousands of files, millions of functions, tens of gigabytes Massive memory/computation complexity for a single machine

WHOPR Architecture - 1





Compilation proceeds in 3 main phases:

- LGEN (Local GENeration)
 - Writes out GIMPLE
 - Produces summary information
- WPA (Whole Program Analysis)
 - Reads summary information
 - Aggregates local callgraphs into global callgraph
 - Produces global optimization plan
- LTRANS (Local TRANSformation)
 - Applies global optimization plan to individual files
 - Performs intra-procedural optimizations
 - Generates final code

WHOPR Architecture - 3



- Phases 1 (LGEN) and 3 (LTRANS) are massively parallel
- Phase 2 (WPA) is fan-in/fan-out serialization point

 Only operates with call graph and symbols
 Transitive closure analysis not computationally expensive
- Scalability provided by splitting analysis and final code generation
 - Restricts types of applicable optimizations
 - For smaller applications, LTRANS provides full IPA functionality (whole program in memory)



Three phases

- Profile code generation: Compile with -fprofile-generate
- Training run: Run code as usual
- Feedback optimization: Recompile with -fprofile-use
- Allows very aggressive optimizations based on accurate cost models
 - Provided that training run is representative!
- Compilation process significantly more expensive
- May not be applicable in all cases



Probes inserted automatically by compiler

Compile and link application with -pg

Run application as usual

Use gprof to analyze output file gmon.out

- \$ gcc -pg -O2 -o matmul matmul.c
- \$./matmul
- \$ gprof ./matmul





- System-wide profiler.
- No modifications to source code
- Samples hardware counters to collect profiling information
- User specifies which hardware counter to sample
- Needs super-user access to start
- Start Oprofiler daemon
- Run application
- Use reporting program to read collected profile

- Instrument \rightarrow Run \rightarrow Recompile cycle too demanding
- New feature being developed to use hardware counters
- 1.Program compiled as usual
- 2.Runs in production environment with hardware counters enabled
- 3.Subsequent recompilations use profile information from hardware counters
- This allows for always-on, transparent profile feedback

GOOg



- There is no "right" choice
- Granularity of work main indicator
- Evaluate complexity ↔ speedup trade-offs
- Combined approach for complex applications
- Algorithms matter!
- Good sequential algorithms may make bad parallel ones

Conclusions

- Performance tuning goes beyond random compiler flags
- Profiling tools are important to study behaviour
- Each tool is best suited for a specific usage
 - Try different flags and use /usr/bin/time to measure
 - Oprofile \rightarrow system wide
 - Gprof \rightarrow intrusive but useful to isolate profiling scope
 - Compiler dumps to determine source of problem
- New advances in instrumentation and whole program compilation will simplify things